# Counting lattice paths by the number of crossings, major index, and descents 

Sergi Elizalde<br>Dartmouth College

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## Overview

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- The number of descents (valleys) and the major index, which were introduced by MacMahon over 100 years ago, and studied by many authors (e.g. yesterday's talk by Terry Visentin).


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- The number of descents (valleys) and the major index, which were introduced by MacMahon over 100 years ago, and studied by many authors (e.g. yesterday's talk by Terry Visentin).

In this talk we will see that combining these statistics one gets surprisingly simple formulas.

## I. Paths crossing a line

## Lattice paths, descents and major index

Let $\mathcal{G}_{a, b}$ be the set of lattice paths in $\mathbb{Z}^{2}$ with a steps $U=(1,1)$ and $b$ steps $D=(1,-1)$, starting at the origin.


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Paths $P \in \mathcal{G}_{a, b}$ can be encoded as binary words via $U \mapsto 0, D \mapsto 1$.

## Definition

- A descent of $P$ is a valley $D U$, let $\operatorname{des}(P)=\#$ descents.


$$
\operatorname{des}(P)=4
$$

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## Definition

- A descent of $P$ is a valley $D U$, let $\operatorname{des}(P)=\#$ descents.
- The major index, maj $(P)$, is the sum of the $x$-coordinates of the descents.



## Lattice paths and major index

$q$-binomial coefficients:

$$
\left[\begin{array}{l}
n \\
k
\end{array}\right]_{q}=\frac{\left(1-q^{n}\right)\left(1-q^{n-1}\right) \cdots\left(1-q^{n-k+1}\right)}{\left(1-q^{k}\right)\left(1-q^{k-1}\right) \cdots(1-q)}
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Lemma (MacMahon)

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\sum_{P \in \mathcal{G}_{a, b}} q^{\operatorname{maj}(P)}=\left[\begin{array}{c}
a+b \\
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a
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$$

## Example

$$
\sum_{P \in \mathcal{G}_{3,2}} q^{\operatorname{maj}(P)}=\left[\begin{array}{l}
5 \\
3
\end{array}\right]_{q}=1+q+2 q^{2}+2 q^{3}+2 q^{4}+q^{5}+q^{6}
$$

## Refinement by the number of descents

Lemma (Fürlinger-Hofbauer '85)

$$
\sum_{P \in \mathcal{G}_{a, b}} t^{\operatorname{des}(P)} q^{\operatorname{maj}(P)}=\sum_{n \geq 0} t^{n} q^{n^{2}}\left[\begin{array}{l}
a \\
n
\end{array}\right]_{q}\left[\begin{array}{l}
b \\
n
\end{array}\right]_{q}
$$

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$$

## Example

$$
\sum_{P \in \mathcal{G}_{3,2}} t^{\operatorname{des}(P)} q^{\operatorname{maj}(P)}=1+t q+2 t q^{2}+2 t q^{3}+\left(t+t^{2}\right) q^{4}+t^{2} q^{5}+t^{2} q^{6}
$$

## Crossing a line

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For $\ell \in \mathbb{Z}$ and $r \geq 0$, let $\mathcal{G}_{a, b}^{\geq r, \ell}$ be the set of paths in $\mathcal{G}_{a, b}$ that cross the line $y=\ell$ at least $r$ times.


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G_{a, b}^{\geq r, \ell}(t, q)=\sum_{P \in \mathcal{G}_{a, b}^{\geq r, \ell}} t^{\operatorname{des}(P)} q^{\operatorname{maj}(P)} .
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For this talk, we focus on the specialization $G_{a, b}^{\geq r, \ell}(q):=G_{a, b}^{\geq r, \ell}(1, q)$.

## Crossing the $x$-axis

First consider the case $\ell=0$, which counts crossings of the $x$-axis.

## Theorem

For any $a, b, r \geq 0$,

$$
G_{a, b}^{\geq r, 0}(q)= \begin{cases}q^{\binom{r+1}{2}}\left[\begin{array}{c}
a+b \\
a+r
\end{array}\right]_{q} & \text { if } a>b, \\
\left(1+q^{a}\right) q^{\binom{r+1}{2}}\left[\begin{array}{c}
2 a-1 \\
a+r
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We give a bijective proof.

## Connections to other work

- The specialization $t=q=1$ (which ignores des and maj) is due to Engelberg '65 and Sen '65, and has later been rediscovered by other authors.


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- The case $t=1$ and $a>b$ can be shown to be equivalent to $a$ result of Seo-Yee '18 about counting ballot paths with marked returns by a different statistic. Their proof is by induction and does not give a bijection.
- The theorem has applications to the enumeration of partitions $\lambda$ with certain restrictions on the ranks $\lambda_{i}-\lambda_{i}^{\prime}$, studied by Corteel-E.-Savage '22+.


## Crossing an arbitrary horizontal line

## Theorem

Let $a, b, m \geq 0$, and let $\ell \in \mathbb{Z} \backslash\{0\}$. If $0<\ell<a-b$, then

$$
G_{a, b}^{\geq 2 m+1, \ell}(q)=G_{a, b}^{\geq 2 m, \ell}(q)=q^{m(2 m+1+\ell)}\left[\begin{array}{c}
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\end{array} .=a-b,\right. \text { then }
\end{aligned}
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$G_{a, b}^{\geq 2 m, \ell}(q)=q^{m(2 m+1+\ell)}\left[\begin{array}{c}a+b \\ a+2 m\end{array}\right]_{q}, \quad G_{a, b}^{\geq 2 m+1, \ell}(q)=q^{m(2 m+1+\ell)}\left[\begin{array}{c}a+b \\ a+2 m+1\end{array}\right]_{q}$.
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# II. Pairs of paths crossing each other 

## Paths with north and east steps

For $A, B \in \mathbb{Z}^{2}$, let $\mathcal{P}_{A \rightarrow B}$ be the set of lattice paths from $A$ to $B$ with steps $N=(0,1)$ and $E=(1,0)$.

$$
A=(x, y) \curvearrowleft \Gamma^{B=(u, v)}
$$

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A descent of $P \in \mathcal{P}_{A \rightarrow B}$ is a corner $E N, \operatorname{des}(P)=\#$ descents,

$$
\begin{aligned}
& \operatorname{des}(P)=2 \\
& A=(x, y)
\end{aligned} \quad \square \square B=(u, v)
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$$
\begin{gathered}
\operatorname{des}(P)=2 \\
A=(x, y) \longmapsto\left\{\begin{array}{l}
7 \\
\operatorname{maj}(P)=2+7, v) \\
2
\end{array}, 9\right.
\end{gathered}
$$

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\square \\
A=(x, y) \longmapsto \\
\operatorname{maj}(P)=2
\end{array}\right)=(u, v) \\
2
\end{gathered}
$$

MacMahon's formula gives

$$
\sum_{P \in \mathcal{P}_{A \rightarrow B}} q^{\operatorname{maj}(P)}=\left[\begin{array}{c}
u-x+v-y \\
u-x
\end{array}\right]_{q}
$$

## Crossings of two paths

A crossing of two paths $P$ and $Q$ is a common vertex $C$ such that:

- $P$ and $Q$ disagree in the step arriving at $C$;
- at the first step after $C$ where $P$ and $Q$ disagree, each path has the same type of step $(N$ or $E$ ) as it had when arriving at $C$.

crossings



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crossings

not a crossing

$$
\begin{aligned}
\mathcal{P}_{\bar{A}_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{\bullet}}^{\geq r}=\{(P, Q): & P \in \mathcal{P}_{A_{1} \rightarrow B_{0}}, Q \in \mathcal{P}_{A_{2} \rightarrow B_{\bullet}} \\
& P \text { and } Q \text { have } \geq r \text { crossings }\} .
\end{aligned}
$$

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A pair in $\mathcal{P}_{\bar{A}_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 3}$ :


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In addition to des and maj, we will keep track of the number of times that the paths cross each other.

For $r \geq 0$, define the polynomials


$$
H_{\bar{A}_{1} \rightarrow B_{\circ}, A_{2} \rightarrow B_{\bullet}}^{\geq r}(t, q)=\sum_{(P, Q) \in \mathcal{P}_{\substack{\geq r}}^{\geq r} B_{\circ}, A_{2} \rightarrow B_{\bullet}} t^{\operatorname{des}(P)+\operatorname{des}(Q)} q^{\operatorname{maj}(P)+\operatorname{maj}(Q)}
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For $r \geq 0$, define the polynomials


$$
\begin{aligned}
& H_{A_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{\mathbf{0}}}^{\geq r}(t, q)=\quad \sum^{\operatorname{des}(P)+\operatorname{des}(Q)} q^{\operatorname{maj}(P)+\operatorname{maj}(Q)} \\
& (P, Q) \in \mathcal{P}_{A_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{0}}^{\geq r}
\end{aligned}
$$

and their specialization $H_{\bar{A}_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{\bullet}}^{\geq r}(q):=H_{\bar{A}_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{\bullet}}^{\stackrel{>}{r}}(1, q)$.

## Easy cases and notation

$$
\text { Let } A_{1}=\left(x_{1}, y_{1}\right), A_{2}=\left(x_{2}, y_{2}\right), B_{1}=\left(u_{1}, v_{1}\right), B_{2}=\left(u_{2}, v_{2}\right) \text {. }
$$

For $r=0$, we can choose the two paths independently, so

$$
H_{A_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{\bullet}}^{\geq 0}(q)=\left[\begin{array}{c}
u_{0}-x_{1}+v_{0}-y_{1} \\
u_{0}-x_{1}
\end{array}\right]_{q}\left[\begin{array}{c}
u_{\bullet}-x_{2}+v_{\bullet}-y_{2} \\
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u_{\bullet}-x_{2}
\end{array}\right]_{q}
$$

To give a general formula, first define
$f_{r, A_{1}, A_{2}, B_{2}, B_{1}}(q):=q^{r\left(r+x_{2}-x_{\mathbf{1}}\right)}\left[\begin{array}{c}u_{2}-x_{1}+v_{2}-y_{1} \\ u_{2}-x_{1}+r\end{array}\right]_{q}\left[\begin{array}{c}u_{1}-x_{2}+v_{1}-y_{2} \\ u_{1}-x_{2}-r\end{array}\right]_{q}$.

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Write $A_{1} \prec A_{2}$ to mean that $A_{1}$ is strictly northwest of $A_{2}$.

## Counting pairs of paths by crossings

## Theorem <br> Let $A_{1}=\left(x_{1}, y_{1}\right), A_{2}=\left(x_{2}, y_{2}\right), B_{1}=\left(u_{1}, v_{1}\right), B_{2}=\left(u_{2}, v_{2}\right)$, where $A_{1} \prec A_{2}$ and $B_{1} \prec B_{2}$, and $x_{1}+y_{1}=x_{2}+y_{2}$.

- $B_{1}$
- $B_{2}$
$A_{1}$ 。
$A_{2}$.


## Counting pairs of paths by crossings

## Theorem

Let $A_{1}=\left(x_{1}, y_{1}\right), A_{2}=\left(x_{2}, y_{2}\right), B_{1}=\left(u_{1}, v_{1}\right), B_{2}=\left(u_{2}, v_{2}\right)$, where $A_{1} \prec A_{2}$ and $B_{1} \prec B_{2}$, and $x_{1}+y_{1}=x_{2}+y_{2}$. Then, for all $m \geq 0$,

$$
H_{A_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 2 m+1}(q)=H_{A_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 2 m}(q)=f_{2 m, A_{1}, A_{2}, B_{2}, B_{1}}(q)
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## Counting pairs of paths by crossings

## Theorem

Let $A_{1}=\left(x_{1}, y_{1}\right), A_{2}=\left(x_{2}, y_{2}\right), B_{1}=\left(u_{1}, v_{1}\right), B_{2}=\left(u_{2}, v_{2}\right)$, where $A_{1} \prec A_{2}$ and $B_{1} \prec B_{2}$, and $x_{1}+y_{1}=x_{2}+y_{2}$. Then, for all $m \geq 0$,

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and for all $m \geq 1$,

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Now let $A=(x, y)$ and $B=(u, v)$. Then, for all $r \geq 0$,

$$
\begin{aligned}
& H_{A \rightarrow B_{1}, A \rightarrow B_{2}}^{\geq}(q)=f_{r, A, A, B_{2}, B_{1}}(q), \\
& H_{A_{1} \rightarrow B, A_{2} \rightarrow B}^{\geq r}(q)=f_{r, A_{1}, A_{2}, B, B}(q),
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$$
\begin{aligned}
H_{A \rightarrow B_{1}, A \rightarrow B_{2}}^{\geq r}(q) & =f_{r, A, A, B_{2}, B_{1}}(q), \\
H_{A_{1} \rightarrow B, A_{2} \rightarrow B}^{\geq r}(q) & =f_{r, A_{1}, A_{2}, B, B}(q), \\
H_{A \rightarrow B, A \rightarrow B}^{\geq r}(q) & = \begin{cases}f_{0, A, A, B, B}(q) & \text { if } r=0, \\
2 \sum_{j \geq 1}(-1)^{j-1} f_{r+j, A, A, B, B}(q) & \text { if } r \geq 1 .\end{cases}
\end{aligned}
$$

## Counting pairs of paths by crossings

With the specialization $t=q=1$ (which ignores des and maj), the theorem still holds without the requirement $x_{1}+y_{1}=x_{2}+y_{2}$.

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For the refinement by maj, we give bijections that have simple descriptions in terms of paths.

For the further refinement by des, our proof is inspired by Krattenthaler's ' 95 refinements of the LGV determinant by des and maj. It is still bijective but relies on certain two-rowed arrays.

# III. Some bijections used in the proofs 

## The bijections $\bar{\tau}$ and $\bar{\sigma}$

Partition $\mathcal{P}_{A \rightarrow B}=\mathcal{P}_{A \rightarrow B}^{E} \cup \mathcal{P}_{A \rightarrow B}^{N}$ according to the last step of the path. Let $\mathbf{v}=(1,-1)$.

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Define a bijection

$$
\bar{\tau}: \mathcal{P}_{A \rightarrow B}^{E} \rightarrow \mathcal{P}_{A+v \rightarrow B}^{N}
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by placing the $N E$ corners of $\bar{\tau}(P)$ at the coordinates of the $E N$ corners of $P$ :


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If $A=(x, y)$ and $B=(u, v)$, one can show that

$$
\operatorname{maj}(\bar{\tau}(P))=\operatorname{maj}(P)+u-x-1
$$

## The bijections $\bar{\tau}$ and $\bar{\sigma}$

Its inverse is the bijection

$$
\bar{\sigma}: \mathcal{P}_{A \rightarrow B}^{N} \rightarrow \mathcal{P}_{A-v \rightarrow B}^{E}
$$

obtained by placing the $E N$ corners of $\bar{\sigma}(P)$ at the coordinates of the NE corners of $P$ :


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obtained by placing the $E N$ corners of $\bar{\sigma}(P)$ at the coordinates of the NE corners of $P$ :


If $A=(x, y)$ and $B=(u, v)$, then

$$
\operatorname{maj}(\bar{\sigma}(P))=\operatorname{maj}(P)-u+x
$$

## A bijection for pairs of paths

Given $(P, Q) \in \mathcal{P}_{\bar{A}_{1} \rightarrow B_{0}, A_{2} \rightarrow B_{0}}^{\geq r}$, let $C$ be the $r$ th crossing from the right. Suppose that $P$ arrives to $C$ with an $N$, and $Q$ with an $E$.


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$$
A_{2}+\mathbf{v}
$$

## A bijection for pairs of paths



With the right setup, this map $(P, Q) \mapsto\left(P^{\prime}, Q^{\prime}\right)$ is a bijection, which we denote by $\theta_{r}$.

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If $A_{1}=\left(x_{1}, y_{1}\right)$ and $A_{2}=\left(x_{2}, y_{2}\right)$, one can show that

$$
\operatorname{maj}\left(P^{\prime}\right)+\operatorname{maj}\left(Q^{\prime}\right)=\operatorname{maj}(P)+\operatorname{maj}(Q)-\left(x_{2}-x_{1}+1\right)
$$

## Composing bijections

To prove our theorem about pairs of paths, we use compositions such as $\theta_{1} \circ \theta_{2} \circ \cdots \circ \theta_{r}$, which decreases maj by $r\left(r+x_{2}-x_{1}\right)$.

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In this example, we have a bijection

$$
\theta_{1} \circ \theta_{2}: \mathcal{P}_{\bar{A}_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 2} \rightarrow \mathcal{P}_{\bar{A}_{1}-2 \mathbf{v} \rightarrow B_{2}, A_{2}+2 \mathbf{v} \rightarrow B_{1}}^{\geq 0} .
$$

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decreases maj by $2\left(2+x_{2}-x_{1}\right)$.
The pairs of paths in the image are easy to enumerate.

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The pairs of paths in the image are easy to enumerate. In this case, with the assumption $A_{1} \prec A_{2}$ and $B_{1} \prec B_{2}$, we obtain

$$
H_{A_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 2}(q)=q^{2\left(2+x_{2}-x_{1}\right)}\left[\begin{array}{c}
u_{2}-x_{1}+v_{2}-y_{1} \\
u_{2}-x_{1}+2
\end{array}\right]_{q}\left[\begin{array}{c}
u_{1}-x_{2}+v_{1}-y_{2} \\
u_{1}-x_{2}-2
\end{array}\right]_{q}
$$

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$H_{A_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{>2}(q)=q^{2\left(2+x_{2}-x_{1}\right)}\left[\begin{array}{c}u_{2}-x_{1}+v_{2}-y_{1} \\ u_{2}-x_{1}+2\end{array}\right]_{q}\left[\begin{array}{c}u_{1}-x_{2}+v_{1}-y_{2} \\ u_{1}-x_{2}-2\end{array}\right]_{q}$.

As another application of these bijections, they can be used to give a simpler proof (directly in terms of paths) of Krattenthaler's refinement by maj of the Lindström-Gessel-Viennot determinantal formula for tuples of non-intersecting paths.

## Bijections for paths crossing a horizontal line

For the problem of a single path crossing a horizontal line, we define similar bijections $\tau$ and $\sigma$. These apply to paths with $U$ and $D$ steps ending on the $x$-axis, and they fix the right endpoint.

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$\tau$ reflects the valleys along the $x$-axis: $\sigma$ reflects the peaks:



$$
\operatorname{maj}(\tau(P))=\operatorname{maj}(P), \quad \operatorname{maj}(\sigma(P))=\operatorname{maj}(P)+\# U-\# D-1
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In this case, we get a bijection $\mathcal{G}_{a, b}^{\geq 2, \ell} \rightarrow \mathcal{G}_{a+2, b-2}$ that decreases maj by $\ell+3$. The paths in the image are easy to count.

## Refinement by the number of descents

Our theorems have refinements that also keep track of des (i.e., the number of $D U$ or $E N$ corners).

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\text { If } 0<\ell<a-b \text {, then }
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$$
G_{a, b}^{\geq 2 m, \ell}(t, q)=\sum_{n} t^{n} q^{n^{2}+m(m+\ell+1)}\left[\begin{array}{c}
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$$

If $A_{1} \prec A_{2}, B_{1} \prec B_{2}$, and $x_{1}+y_{1}=x_{2}+y_{2}$, then, for all $m \geq 0$,

$$
\left.\begin{array}{rl}
H_{A_{1} \rightarrow B_{2}, A_{2} \rightarrow B_{1}}^{\geq 2 m}(t, q) \\
= & q^{2 m\left(2 m+x_{2}-x_{1}\right)}
\end{array}\right)\left(\sum_{n} t^{n} q^{n(n+2 m)}\left[\begin{array}{c}
u_{2}-x_{1} \\
n
\end{array}\right]_{q}\left[\begin{array}{c}
v_{2}-y_{1} \\
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## THANK YOU

References:

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