

MATH 53

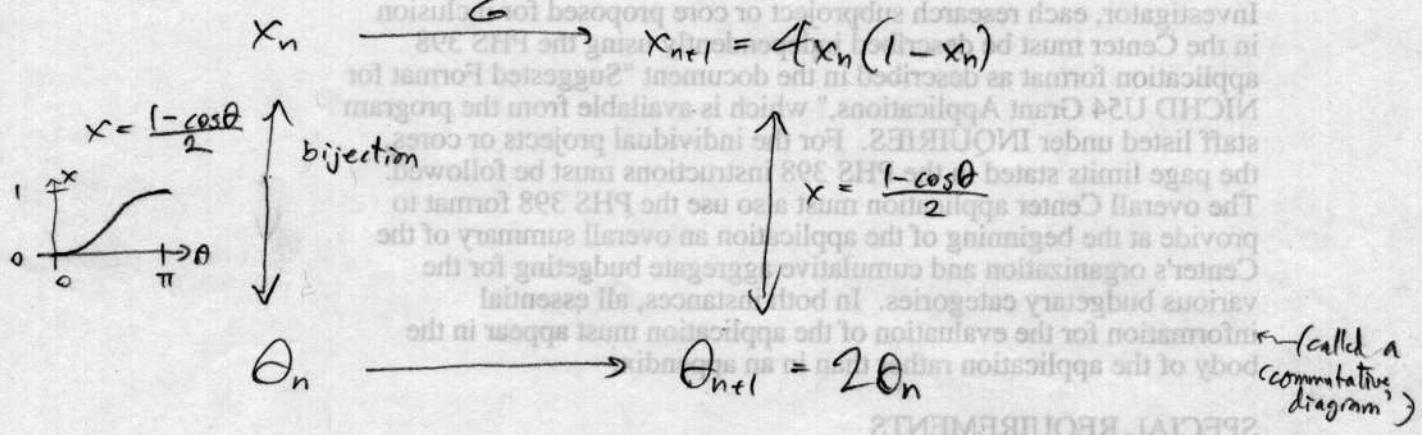
: proof of exponential smallness of  
itinerary subintervals.

Banoff  
10/1/69

[This will also explain 1.15 from HW1 ... the trig one!]

It turns out if you map  $x$  to a new variable  $\theta$ , then the logistic map  $G(x) = 4x(1-x)$  becomes ridiculously simple: just doubling  $\theta$ .

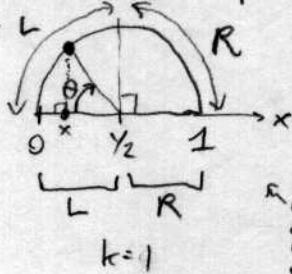
I.e:



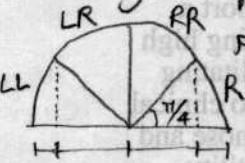
this states that iterating  $x_n$  to  $x_{n+1}$  via  $G$  is equivalent to converting to  $\theta$ , doubling  $\theta$ , then converting back again. (repeat this giving  $\theta \rightarrow 2^k \theta$  for HW1 E.15)

Why does this work? If  $x = \frac{1-\cos\theta}{2}$  then  $\frac{1-\cos 2\theta}{2} = \sin^2 \theta = 1 - \cos^2 \theta = 4\left(\frac{1-\cos\theta}{2}\right)\left(\frac{1+\cos\theta}{2}\right)$

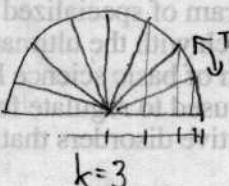
Itineraries of the doubling map  $\theta \rightarrow 2\theta$  are really easy to compute:



note how  
geometrically,  
 $x = \frac{1-\cos\theta}{2}$



$k=2$



$k=3$



Projecting the semicircle down to  $(0, 1]$ , we see that no subinterval of  $k$  symbols may exceed  $\frac{1}{2} \cdot \frac{\pi}{2^k} = \frac{\pi}{2^{k+1}}$  in length along  $x$ . QED.

the subintervals are equally spaced in angle since doubling is simple