

Nonuniform FFTs at Flatiron $-- \sim * \quad OR \quad * \sim --$

Lessons from developing a small numerical library

Alex Barnett¹, with much help from: Jeremy Magland¹, Ludvig af Klinteberg (Mälardalen U.), Melody Shih (NVidia), Joakim Andén (KTH), Libin Lu¹, Robert Blackwell (SCC), Andrea Malleo (Bloomberg), and many others...

FWAM5 Friday, October 20, 2023

¹Center for Computational Mathematics, Flatiron Institute





Given locations x_1, \ldots, x_M , and their "strengths" c_1, \ldots, c_M , return

$$f_k = \sum_{j=1}^M e^{ikx_j} c_j$$
, for $k = -\frac{N}{2}, -\frac{N}{2} + 1, \dots, \frac{N}{2} - 2, \frac{N}{2} - 1$

- Has N outputs, each a sum of M terms: naively $\mathcal{O}(NM)$ cost (flops)
- Off-grid version of fast Fourier transform (FFT) there M = N and $x_j = 2\pi j/N$





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"Type 3" as type 1 but arbitrary target freqs. $\{s_k\}_{k=1}^N$: $f_k = \sum_{j=1}^M e^{is_k imes_j} c_j$ $NU \rightarrow NU$

Higher dimensions also needed 2D type 1: $f_{k,\ell} = \sum_{i=1}^{M} e^{i(kx_j + \ell y_i)} c_j$ for k, ℓ in a rectangle of modes Ħ regular good (U) nonuniform (NU) pts. $f_{k,\ell,m} = \sum_{i=1}^{M} e^{i(kx_j + \ell y_j + mz_j)} c_i$

3D type 1: $f_{k,\ell,m} = \sum_{j=1}^{m} e^{i(kx_j + \ell y_j + mz_j)} c_j$ etc

- dimensions {1,2,3} × types {1,2,3} = 9 transforms
 Software design: how to avoid code repetition?
- 9 transforms $~\times~~\{{\sf float},{\sf double}\} = 18$ functions

How reduce number of functions to write and maintain?



Who uses such transforms?

- 1) Fourier image reconstruction: $f_{k,\ell}$ is unknown pixel intensities apparatus measures strengths c_i at nonuniform frequency points (x_i, y_i)
 - MRI (either 2D slice, or 3D)
 - coherent diffraction/powder imaging (X-ray)
 - very long baseline interferometry (VLBI)
 - cryo electron microscopy



some 2D MRI points



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- 2) Numerical forward solvers, simulation
 - electrostatics or fluid problems in periodic box

spectral Ewald method: Poisson solve trivial in Fourier space

- numerical PDE eg, interpolating between overlapping grids
- eval Fourier *transform* by numerical quadrature (type 1)
- Fresnel diffraction (optics)



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- 3) Spatial/temporal statistics
 - power-spectrum of NU time-series, or point-masses galaxies
 - fast kernel apply in Gaussian process regression



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Probably the most famous astro image of 2019:

10⁻¹⁰ radian resolution!



uniform image grid $f_{k,\ell}$





Telescope Locations

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predicted signals $\{g_i\}$ at NU pts

• our library (FINUFFT) used by SMILI code for this image (K. Akiyama, '19)

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predicted signals $\{g_i\}$ at NU pts

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How? the above is a linear *forward model*: g = Af big dense matrix A lterative optimization of f until it best fits the detected signal:

$$\mathbf{f}_{\text{recon}} = \arg\min_{\mathbf{f}} \|A\mathbf{f} - \mathbf{g}_{\text{detected}}\|_2^2 + \lambda_1 \|\mathbf{f}\|_1 + \lambda_{\text{TV}} \|\mathbf{f}\|_{\text{TV}}$$

• each iteration, A and A^* applied fast by NUFFTs dominant cost, I think

same idea in other 2D or 3D Fourier imaging (MRI, cryo-EM, etc)

(Dutt-Rokhlin '93, Steidl '98, Greengard-Lee '04, Potts et al, ...)

set up new grid on $[0, 2\pi)$ with n = 2N points, say

Eg, 1D Type 1:

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Design a kernel $\psi(x)$

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fun math: small width (w grid-points) YET ε -small Fourier tails

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1) Spread each spike c_j onto grid $b_\ell = \sum_{j=1}^M c_j \psi(\ell h - x_j)$ detail: 2π -periodize



 $\mathcal{O}(w^d M)$ flops $\mathcal{O}(wdM)$ kernel evals often dominates cost

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keep only low modes $-N/2 \le k < N/2$

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Type 2 reverses the steps; Type 3 is "Type 2 wrapped inside a Type 1"

Prehistory (2015): Leslie Greengard, Jeremy Magland,



Marina Spivak, myself, at SCDA NUFFT needed for imaging (cryo-EM, etc) had: NYU single-threaded Fortran, Gaussian kernel $\psi(x)$ too wide (Greengard-Lee '09) also: NFFT3 C++, multithreaded, hard to use, user chooses $\psi(x)$ (Keiner et al '06) Wanted faster multithreaded code, easy-to-use from many languages



finufft

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Performance: kernel $\psi(x)$ evaluation

Problem: for some chips & compilers, exp(x) slow (40 M evals/sec/core) Fix: piecewise polynomial approx + Horner's rule C gen. by MATLAB!



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```
/ Code generated by gen_all_horner_C_code.m in finufft/devel
// Authors: Alex Barnett & Ludvig af Klinteberg.
// (C) The Simons Foundation. Inc.
  if (w==2) {
   FLT c0[] = {4.5147043243215315E+01, 4.5147043243215300E+01, 0.00000000000000E+00, 0.00000000000000E+00};
   FLT c3[] = {-2.0382426253182082E+01, 2.0382426253182086E+01, 0.00000000000000E+00, 0.00000000000000E+00};
   FLT c4[] = {-2.0940804433577420E+00, -2.0940804433577389E+00, 0.00000000000000E+00, 0.00000000000000E+00};
   for (int i=0: i<4: i++) ker[i] = c0[i] + z^{(c1[i])} + z^{(c2[i])} + z^{(c3[i])} + z^{(c4[i])));
  } else if (w==3) {
   FLT c1[] = {3.1653018869611077E+02, 7.4325702843759617E-14, -3.1653018868907071E+02, 0.0000000000000000E+00};
   FLT c2[] = \{1,7742692790454484E+02, -3,3149255274727801E+02, 1,7742692791117119E+02, 0,000000000000000E+00\}
   FLT c5[] = {-3,9654011076088804E+00, 1,8062124448285358E-13, 3,9654011139270540E+00, 0,000000000000000000E+00};
   for (int i=0; i<4; i++) ker[i] = c0[i] + z*(c1[i] + z*(c2[i] + z*(c3[i] + z*(c4[i] + z*(c5[i])))));</pre>
  } else if (w==4) {
   FLT c0[] = {5,4284366850213200E+02, 1,0073871433088398E+04, 1,0073871433088396E+04, 5,4284366850213223E+02};
   FLT c1[] = {1,4650917259256939E+03, 6,1905285583602863E+03, -6,1905285583602881E+03, -1,4650917259256937E+03};
   FLT c2[] = {1.4186910680718345E+03, -1.3995339862725591E+03, -1.3995339862725598E+03, 1.4186910680718347E+03};
   FLT c3[] = {5.1133995502497419E+02, -1.4191608683682996E+03, 1.4191608683682998E+03, -5.1133995502497424E+02];
   FLT c4[] = {-4.8293622641174039E+01, 3.9393732546135226E+01, 3.9393732546135816E+01, -4.8293622641174061E+01};
   FLT c5[] = {-7.8386867802392288E+01, 1.4918904800408930E+02, -1.4918904800408751E+02, 7.8386867802392359E+01};
   FLT c6[] = {-1.0039212571700894E+01, 5.0626747735616746E+00, 5.0626747735625512E+00, -1.0039212571700640E+01};
   for (int i=0; i<4; i++) ker[i] = c0[i] + z*(c1[i] + z*(c2[i] + z*(c3[i] + z*(c4[i] + z*(c5[i] + z*(c6[i]))))));
  } else if (w==5) {
   FLT c0[] = {9.9223677575398392E+02, 3.7794697666613320E+04, 9.8715771010760494E+04, 3.7794697666613283E+04, 9.9223
$677575398403E+02, 0.0000000000000000E+00, 0.000000000000000E+00, 0.00000000000000E+00};
-:--- ker horner allw loop.c Top L1 Git-master (C/*L WS Abbrev)
```

- GCC/ICC compilers SIMD-vectorize this; get 400-700 M evals/sec/core
- think hard re SIMD, but avoid maintaining intrinsics immintrin.h
- are exploring custom AVX512 (Wenda Zhou + R. Blackwell)



Performance: spreading

The order in which NU points spread to grid has big effect on speed!

- bin-sort NU pts (into $16 \times 4 \times 4$ cuboids of grid)
- process all pts in bin 1, then bin 2, ... good for keeping grid in cache(J. Magland)
- multithreaded the bin-sort

(2023: M. Reinecke speeds it up!)

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Load-balanced multithreading

Type-1 not so: writes collide!

(also collab with J. Magland)

Type-2 easy: OpenMP parallel over NU pts no collisons reading from U blocks load-balance via "subproblems" each of 10⁴ NU pts



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C-compatible: pass pointers, explicit array sizes, return value is error code...



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int M = 1e8;  // number of nonuniform points
vector<double> x(M);  // NU pts
vector<complex<double> > c(M);  // NU strengths
// (here user fills x and c as they like)
int N = 1e7;  // number of output modes
vector<complex<double> > f(N);  // allocate output array
int status = finufft1d1(M, &x[0], &c[0], +1, 1e-9, N, &f[0], NULL);  // do it
```

• Why last arg NULL? it accepts ptr to options C-struct eg, opts.debug=2



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2020: maintaining 18 × 2 functions too much pain → "guru" interface
4-function pattern: Create plan, Set the NU pts, Execute transform(s), Destroy plan.
finufft_plan: "opaque" pointer to (private) C++ struct. (as in Brian's talk)



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finufft_plan: "opaque" pointer to (private) C++ struct. (as in Brian's talk)
Lesson: learn about public vs private headers, namespacing (as I had to)
Lesson: think hard about interface, break it VERY rarely!
help users: preserve all simple and batch interfaces (they call guru)

Wrappers to other languages: expands user base

Guru interface made wrapping easier: just wrap 4 funcs, pass opaque ptr write the simple and batched functions via a few lines *in each high-level language*

MATLAB/Octave

(Libin Lu)

(Bindel '09)

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still recommend MWrap to auto-gen MEX code horror



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served well for a while. But eg forced recompile of libfinufft.so

2020: Joakim Andén switched us to ctypes nice, 1000 lines incl. auto-doc-gen

Others have wrapped (cu)FINUFFT in autodiff frameworks: tensorflow-nufft, jax-finufft, pytorch-finufft



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2018: Dan Foreman-Mackey wrote a pybind11 wrapper to FINUFFT

served well for a while. But eg forced recompile of libfinufft.so

2020: Joakim Andén switched us to ctypes nice, 1000 lines incl. auto-doc-gen

Others have wrapped (cu)FINUFFT in autodiff frameworks: tensorflow-nufft, jax-finufft, pytorch-finufft

Julia wrapper is separate repo: helps separate concerns (L. af Klinteberg) Lessons: each new language brings installation troubles linux/OSX/Windows

Need tests that check accuracy for all transforms fail if measured error $\geq 10\varepsilon$?

> make test [includes bunch of edge cases M=0, N=0, eps=0.0, etc...]
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A story: I wrote finufft1d_test, etc, writes to stdout for humans...

```
test 1d type 1:
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[...]
Such speed and accuracy testers are crucial for progress
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Documentation

Next topic

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README.md, or a .tex file to PDF, fine for small project, but:



FINUFFT is a library to compute efficiently the three most common types of nonuniform fast Fourier transform (NUFFT) to a specified precision, in one, two, or three dimensions, either on a multi-core shared-memory machine, or on a GPU. It is extremely fast (typically activity of 10⁴ to 10⁵ points per second on a GPU, or up to 10⁶ points per second on a GPU, has very simple interfaces to most major numerical languages (CC++, Fortan, MATLB), cotxey, Python, and Julia), but also has more advanced (vectorized and "guru") interfaces that allow multiple strength vectors and the rouse of FFTN. It has been developed since 2017 at the Center for Computational Mathematics at the Flatton Institute, by Alex Barnett and others, and is released under an Apple v2 (lonse).

What does FINUFFT do?

As an example, given M real numbers $x_j \in [0, 2\pi)$, and complex numbers c_j , with $j = 1, \ldots, M$, and a requested integer number of modes N, FINUFFT can efficiently compute the 1D "type 1" transform, which means to evaluate the N complex outputs

$$f_k = \sum_{i=1}^{M} c_j e^{ikx_j}$$
, for $k \in \mathbb{Z}$, $-N/2 \le k \le N/2 - 1$. (1)

As with other "fast" algorithms, FINUFFT does not evaluate this sum directly—which would take O(NM) effort—but rather uses a sequence of stags (in this case, optimaly chosen sprateding, FFT, and deconvolution) to approximate the vector of answers (1) to within the user's desired relative toler-ace, with only $O(N\log N + AI)$ efforts. (in, almost timear time. Thus the speed-up is similar to that of the FFT. You may now want to jump to quickstart, or see the definitions of the other transforms in general dimension.

One interpretation of (1) is: the returned values f_k are the Fourier series coefficients of the fourier series is distribution $f(x) := \sum_{j=1}^{M} c_j \delta(x - x_j)$, a sum of point-masses with arbitrary locations x_j and strengths c_i . Such exponential sums are needed in many applications in science and engineering, in-

Web-facing essential for professionalism

- math definitions
- installation
- examples
- tutorial cases
- troubleshooting
- devnotes



 Decent solution I recommend: sphinx, host at readthedocs.io docs/*.rst ReStructuredText, in master branch (unlike GHpages)
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But don't get carried away with obscure doc build automation eg: my bash script auto-generates C/C++ docs, now hard to understand



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- Always update docs at the time you change the source



GPU version

cuFINUFFT: CUDA called from C++ (Melody Shih, intern '18, '19 [NYU \rightarrow NVidia])



 spreading done in shared memory (subproblems), not global memory (Shih et al, PDSEC 2021, best paper prize)

2023: cleaned up GPU lib (+tests,...) into FINUFFT repo (R. Blackwell)

Future: language interfaces auto-detect GPU arrays?



First identify a need

+ benchmark existing codes, on same task!

- can you perform task faster / more accurate?



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Start small, but professionalizing a library involves learning!

- users find bugs you'd never thought of must respond (thanks Libin)
- users make improvement PRs, or beat your speed respond to some
- be conservative: one change at a time for your sanity (eg: no MPI)
- try to be somewhat idiomatic $\,$ I won't use #define FLT double in C++ again \odot



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Installation, dependencies & language wrappers takes much of our time!

- helps users, but needs a team (linux,OSX,Windows). Your choice.

Eg: guess which language wrapper SMILI black hole astro code uses?

- still plenty of fun with math & floating-point speed hacking!

Thank-you!

